

Constraint Systems Cheat Sheet

	Boolean Algebra B	Rational Trees RT	Linear Polynomial Equations \mathfrak{R}	Finite Domains FD																																																		
Domain	Truth values 0 and 1	Herbrand universe with Herbrand interpretation	The set \mathfrak{R} of real numbers	The set \mathcal{Z} of integers																																																		
Allowed atomic constraints	$C ::= X = Y \mid \neg X = Y \mid X \odot Y = Z$, where X, Y and Z are variables or truth values.	$C ::= s \doteq t$ (s, t : terms over Σ)	$C ::= a_1 * X_1 + \dots + a_n * X_n + b \odot 0$, ($n \geq 0$) <ul style="list-style-type: none"> coefficients $a_i, b \in \mathfrak{R}$, $a_i \neq 0$, variables X_1, \dots, X_n totally ordered in strictly descending order $\odot \in \{=, <, \leq, >, \geq, \neq\}$ 	$C ::= X \text{ in } n..m \mid X \text{ in } [k_1, \dots, k_l] \mid$ $X \odot Y \mid X + Y = Z$ <ul style="list-style-type: none"> n, m, k_1, \dots, k_l: integers ($l \geq 0$) $\odot \in \{=, <, >, \leq, \geq, \neq\}$ X, Y and Z: pairwise distinct variables 																																																		
Constraint Theory	<p>Instances of $\neg X = Z$ and $X \odot Y = Z$ according to the following truth table, where $\odot \in \{\sqcap, \sqcup, \oplus, \rightarrow, \leftrightarrow\}$.</p> <table border="1" style="display: inline-table; border-collapse: collapse; text-align: center;"> <thead> <tr> <th>X</th> <th>Y</th> <th>$\neg X$</th> <th>$X \sqcap Y$</th> <th>$X \sqcup Y$</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>1</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>0</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>0</td><td>1</td><td>1</td></tr> </tbody> </table> <table border="1" style="display: inline-table; border-collapse: collapse; text-align: center;"> <thead> <tr> <th>X</th> <th>Y</th> <th>$X \oplus Y$</th> <th>$X \rightarrow Y$</th> <th>$X \leftrightarrow Y$</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td><td>1</td><td>1</td></tr> <tr><td>0</td><td>1</td><td>1</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>0</td><td>1</td><td>1</td></tr> </tbody> </table>	X	Y	$\neg X$	$X \sqcap Y$	$X \sqcup Y$	0	0	1	0	0	0	1	1	0	1	1	0	0	0	1	1	1	0	1	1	X	Y	$X \oplus Y$	$X \rightarrow Y$	$X \leftrightarrow Y$	0	0	0	1	1	0	1	1	1	0	1	0	1	0	0	1	1	0	1	1	<p>Clark's Equality Theory Without Acyclicity</p> <p>Reflexivity $\forall (true \rightarrow x \doteq x)$</p> <p>Symmetry $\forall (x \doteq y \rightarrow y \doteq x)$</p> <p>Transitivity $\forall (x \doteq y \wedge y \doteq z \rightarrow x \doteq z)$</p> <p>Compatibility $\forall (x_1 \doteq y_1 \wedge \dots \wedge x_n \doteq y_n \rightarrow f(x_1, \dots, x_n) \doteq f(y_1, \dots, y_n))$</p> <p>Decomposition $\forall (f(x_1, \dots, x_n) \doteq f(y_1, \dots, y_n) \rightarrow x_1 \doteq y_1 \wedge \dots \wedge x_n \doteq y_n)$</p> <p>Contradiction (Clash) $\forall (f(x_1, \dots, x_n) \doteq g(y_1, \dots, y_m) \rightarrow false)$ if $f \neq g$ or $n \neq m$</p>	<p>Tarski's theory of real closed fields</p> $C_1 \quad (x + y) + z = x + (y + z)$ $C_2 \quad x + 0 = x$ $C_3 \quad x + (-1 * x) = 0$ $C_4 \quad x + y = y + x$ $C_5 \quad (x * y) * z = x * (y * z)$ $C_6 \quad x * 1 = x$ $C_7 \quad x \neq 0 \rightarrow \exists y \ x * y = 1$ $C_8 \quad x * y = y * x$ $C_9 \quad x * (y + z) = (x * y) + (x * z)$ $C_{10} \quad 0 \neq 1$ $O_1 \quad \neg(x < x)$ $O_2 \quad x < y \wedge y < z \rightarrow x < z$ $O_3 \quad x < y \vee x = y \vee y < x$ $O_4 \quad x < y \rightarrow x + z < y + z$ $O_5 \quad 0 < x \wedge 0 < y \rightarrow 0 < x * y$ $R_1 \quad 0 < x \rightarrow \exists y \ y * y = x$ $R_2 \quad y_n \neq 0 \rightarrow \exists x \ y_n * x^n + y_{n-1} * x^{n-1} + \dots + y_0 = 0$ for odd n	<p>Presburger Arithmetic</p> $0 = s(X) \rightarrow \perp$ $X = X$ $X = Y \rightarrow s(X) = s(Y)$ $s(X) = s(Y) \rightarrow X = Y$ $X = Y \wedge Y = Z \rightarrow X = Z$ $X + 0 = X$ $X + s(Y) = s(X + Y)$ $X \leq Y \leftrightarrow \exists Z. X + Z = Y$ $X \text{ in } n..m \leftrightarrow n \leq X \wedge X \leq m$ $X \text{ in } [k_1, \dots, k_l] \leftrightarrow X = k_1 \vee \dots \vee X = k_l$
X	Y	$\neg X$	$X \sqcap Y$	$X \sqcup Y$																																																		
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CHR rules	<p>a1 @ and(X,Y,Z) <=> X=0 Z=0. a2 @ and(X,Y,Z) <=> Y=0 Z=0. a3 @ and(X,Y,Z) <=> X=1 Y=Z. a4 @ and(X,Y,Z) <=> Y=1 X=Z. a5 @ and(X,Y,Z) <=> X=Y Y=Z. a6 @ and(X,Y,Z) <=> Z=1 X=1, Y=1.</p> <p>Analogous for other logical connectives.</p>	<p>X eq X <=> var(X) true. T eq X <=> var(X), X << T X eq T. T1 eq T2 <=> nonvar(T1), nonvar(T2) same_functor(T1, T2), args2list(T1, L1), args2list(T2, L2), same_args(L1, L2). X eq T1, X eq T2 <=> var(X), X << T1, T1 = << T2 X eq T1, T1 eq T2.</p>	<p>eliminate @ A1*X+P1 eq 0 \ PX eq 0 <=> find(A2*X, PX, P2) normalize(A2*(-P1/A1)+P2, P3), P3 eq 0. empty @ B eq 0 <=> number(B) zero(B).</p>	<p>X in A..B <=> A>B false. X in A..B, X in C..D <=> X in max(A,C)..min(B,D). X le Y, X in A..B, Y in C..D <=> B>D X le Y, X in A..D, Y in C..D. X le Y, X in A..B, Y in C..D <=> C<A X le Y, X in A..B, Y in A..D. X eq Y, X in A..B, Y in C..D <=> A\=C X eq Y, X in max(A,C)..B, Y in max(C,A)..D. X eq Y, X in A..B, Y in C..D <=> B\D X eq Y, X in A..min(B,D), Y in C..min(D,B). X ne Y, X in A..B, Y in C..D <=> A=C, C=D X ne Y, X in (A+1)..B, Y in C..D.</p>																																																		