

## Theorem Proving - Resolution

**Boolean CSP in CNF:** Conjunction of clauses

**Clause:** Disjunction of Literals

**Literal:** Positive or negative atomic proposition

## High-Level Implementation in CHR

Syntax: Clause as *ordered* list of signed variables

For example,  $\neg x \vee y \vee z$  as `cl([-x,+y,+z])`.

## Resolution with Factoring

```
empty_clause @ cl([]) <=> false.
```

```
tautology @ cl(L) <=> in(+X,L),in(-X,L) | true.
```

```
cl(L1),cl(L2) ==> del(+X,L1,L3),del(-X,L2,L4) |
    merge(L3,L4,L),
    cl(L).
```

Auxiliary predicates:

`in(A,L)`: Element A occurs in list L.

`del(A,L1,L2)`: List L1 without element A is L2.

`merge(L1,L2,L)`: Ordered merge of L1 and L2 is L3.

Sorry, examples are missing. Try it online with webchr.

## Subsumption

Subsumption

$\text{cl}(\text{L1}), \text{cl}(\text{L2}) \Leftrightarrow \text{sublist}(\text{L1}, \text{L2}) \mid \text{cl}(\text{L1}).$

From (Subsumption) and (Resolution)

$\text{cl}(\text{L1}), \text{cl}(\text{L2}) \Rightarrow \text{del}(+X, \text{L1}, \text{L3}), \text{del}(-X, \text{L2}, \text{L4}) \mid$   
 $\text{merge}(\text{L3}, \text{L4}, \text{L}),$   
 $\text{cl}(\text{L}).$

derive Parental Subsumption

$\text{cl}(\text{L1}), \text{cl}(\text{L2}) \Leftrightarrow \text{del}(\text{A}, \text{L1}, \text{L3}), \text{del}(\text{B}, \text{L2}, \text{L4}),$   
 $\text{compl}(\text{A}, \text{B}), \text{sublist}(\text{L3}, \text{L4}) \mid$   
 $\text{cl}(\text{L1}), \text{cl}(\text{L4}).$

Auxiliary predicates:

**sublist(L1,L2)**: All elements of L1 occur in L2.

**compl(A,B)**: Literal A is the logical complement of B.

## Unit Resolution Davis-Putnam Procedure

From (Subsumption)

$$\text{cl}(L1), \text{cl}(L2) \Leftrightarrow \text{sublist}(L1, L2) \mid \text{cl}(L1).$$

Unit Subsumption

$$\text{cl}([A]), \text{cl}(L) \Leftrightarrow \text{in}(A, L) \mid \text{cl}([A]).$$

From (Parental Subsumption)

$$\begin{aligned} \text{cl}(L1), \text{cl}(L2) \Leftrightarrow & \text{del}(A, L1, L3), \text{del}(B, L2, L4), \\ & \text{compl}(A, B), \text{sublist}(L3, L4) \mid \\ & \text{cl}(L1), \text{cl}(L4). \end{aligned}$$

Unit Propagation

$$\begin{aligned} \text{cl}([A]), \text{cl}(L) \Leftrightarrow & \text{compl}(A, B), \text{del}(B, L, L1) \mid \\ & \text{cl}([A]), \text{cl}(L1). \end{aligned}$$

Incomplete without search

## Ordered Resolution

(Resolution)

```
cl(L1),cl(L2) ==> del(+X,L1,L3),del(-X,L2,L4) |
    merge(L3,L4,L),
    cl(L).
```

Ordered Resolution I

```
cl([A|L1]),cl(L2) ==> compl(A,B),del(B,L2,L3) |
    merge(L1,L3,L),
    cl(L).
```

Ordered Resolution II

```
cl([A|L1]),cl([B|L2]) ==> compl(A,B) |
    merge(L1,L2,L),
    cl(L).
```

Still satisfaction complete

## Clause Length Filtering

(Resolution)

```
cl(L1),cl(L2) ==> del(+X,L1,L3),del(-X,L2,L4) |
                      merge(L3,L4,L),
                      cl(L).
```

Filtered Resolution

```
cl(L1),cl(L2) ==> maxlen(L1,n1),maxlen(L2,n2),
                      compl(A,B),
                      del(A,L1,L3),del(B,L2,L4) |
                      merge(L3,L4,L),
                      cl(L).
```

$n_1 = 1, n_2 = \infty$	Unit propagation
$n_1 = 1, n_2 = 1$	Contradiction
$n_1 = 2, n_2 = 2$	Binary Resolution
$n_1 = 2, n_2 = 3$	Ternary Resolution

Incomplete without search

## Search in CHR<sup>∨</sup>

Retain completeness for restricted resolution by  
 Search = Enumeration = Labeling =  
 Branching = Case Analysis = Choices

Variables

```
enum([X|L]) <=> (cl([+X]) ; cl([-X])),  
                      enum(L).
```

Clauses I

```
label, cl(L) <=> in(A,L), cl([A]), label.
```

Clauses II

```
split, cl(L) <=> split(L,L1,L2),  
                           (cl(L1) ; cl(L2)),  
                           split.
```

The semicolon (;) denotes disjunction implemented by search.

## Constraint Inference

Derive equality and partial order constraints.

Equivalence

$\text{cl}([+X, -Y]), \text{cl}([-X, +Y]) \Leftrightarrow X = Y.$

Replace all  $X$  by  $Y$ .

Implication

$\text{cl}([+X, -Y]) \Rightarrow X \geq Y.$

Helps finding equivalences.

Exclusive Or

$\text{cl}([+X, +Y]), \text{cl}([-X, -Y]) \Leftrightarrow X \neq Y.$

Replace all  $X$  by  $-Y$ .

## Projection by Variable Elimination

(Resolution)

```
cl(L1),cl(L2) ==> del(+X,L1,L3),del(-X,L2,L4) |  
                      merge(L3,L4,L),  
                      cl(L).
```

Variable Elimination

```
eliminate(X),  
cl(L1),cl(L2) ==> del(+X,L1,L3),del(-X,L2,L4) |  
                      merge(L3,L4,L),  
                      cl(L).
```

```
eliminate(X),  
cl(L) <=> (in(+X,L);in(-X,L)) | true.
```

```
eliminate(X) <=> true.
```

## Projection by Variable Elimination

Example

```
?- cl([-x,+y]), cl([-x,+z]), cl([+x,-y,-z]),
eliminate(y).
```

```
cl([-x),+(z)])
```

```
?- cl([-x,+y]), cl([-x,+z]), cl([+x,-y,-z]),
eliminate(z).
```

```
cl([-x),+(y)])
```

```
?- cl([-x,+y]), cl([-x,+z]), cl([+x,-y,-z]),
eliminate(x), eliminate(y).
```

true

```
?- cl([-x,+y]), cl([+x]), cl([-y]),
eliminate(x), eliminate(y).
```

false

## Full First Order Resolution

(Resolution)

```
cl(L1),cl(L2) ==> del(+X,L1,L3),del(-X,L2,L4) |
    merge(L3,L4,L),
    cl(L).
```

First Order Resolution

```
cl(C1),cl(C2) ==> copy(C1,L1),copy(C2,L2),
    del(A,L1,L3),del(B,L2,L4),
    compl(A,B),unify(A,B) |
    merge(L3,L4,L),
    cl(L).
```